

## Homework Assignment 2

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(Submission deadline: April 16th 10:20, 2024, Beijing time, before the class starts. **No submission after the deadline will be accepted.**)

**Notations:**  $\mathbb{N}$  denotes the set of natural numbers. For a set  $X$ , we write  $x \in X$  when  $x$  is an element of  $X$ .

**Question 1.** Suppose  $\Sigma$  is a set of formulas that is inconsistent. Prove  $\Sigma \vdash \varphi$  holds for any formula  $\varphi$ .

*Hint: get  $\perp \rightarrow (\neg\varphi \rightarrow \perp)$  and  $(\neg\varphi \rightarrow \perp) \rightarrow \varphi$  from axioms, and then derive  $\varphi$  by Modus Ponens.  $(\neg\varphi \rightarrow \perp) \rightarrow \varphi$  is easy to find, how about  $\perp \rightarrow (\neg\varphi \rightarrow \perp)$ ? Recall  $(\varphi \rightarrow \psi)$  means  $(\neg\varphi \vee \psi)$ .*

**Question 2.** Prove the opposite of Deduction Lemma; that is, if  $\Sigma \vdash (\varphi \rightarrow \psi)$  then  $\Sigma \cup \{\varphi\} \vdash \psi$ .

*Hint: by assumption, there is a proof  $\varphi_1, \dots, \varphi_n$  of  $(\varphi \rightarrow \psi)$  from  $\Sigma$ . How can you modify it into a proof of  $\psi$  from  $\Sigma \cup \{\varphi\}$ ?*

**Question 3.** Suppose a set of formulas  $\Sigma_n$  is given for each  $n \in \mathbb{N}$ , and suppose  $\Sigma_n \subseteq \Sigma_{n+1}$  holds for each  $n$ . Further assume  $\Sigma_n$  is consistent for each  $n$ . Under these assumptions, prove  $\Sigma_\infty = \bigcup_{n \in \mathbb{N}} \Sigma_n$  (i.e., the union of  $\Sigma_n$  for all  $n$ ) is also consistent.

*Hint: prove by contradiction; that is, assume  $\Sigma_\infty$  is inconsistent and derive contradiction. In particular, if  $\Sigma_\infty$  is inconsistent, then you can show  $\Sigma_n$  is inconsistent for some  $n$ .*

**Question 4.** Let  $\Sigma$  be a complete set of formulas (p14 of slides, Lecture 5). Prove  $(\varphi_1 \vee \varphi_2) \in \Sigma$  if and only if  $\varphi_1 \in \Sigma$  or  $\varphi_2 \in \Sigma$ , in the following steps:

1. Prove, for a given set  $\Sigma'$  of formulas and  $\varphi$ , if  $\Sigma' \vdash \varphi$  and  $\Sigma' \vdash \neg\varphi$ , then  $\Sigma'$  is inconsistent.
2. Prove, for any  $\varphi$ , we have  $\Sigma \vdash \varphi$  if and only if  $\varphi \in \Sigma$ .

3. Prove, for any  $\varphi_1, \varphi_2$ , we have  $(\varphi_1 \vee \varphi_2) \in \Sigma$  if and only if  $\varphi_1 \in \Sigma$  or  $\varphi_2 \in \Sigma$ .

*Hint: For 1, find a suitable axiom and derive  $\Sigma' \vdash \perp$ . For 2, it is easy to show  $\varphi \in \Sigma \Rightarrow \Sigma \vdash \varphi$ ; then prove the opposite implication by contradiction, using 1. For 3, assume  $\varphi_1 \in \Sigma$  or  $\varphi_2 \in \Sigma$  and show  $\Sigma \vdash (\varphi_1 \vee \varphi_2)$ , then use 2; this proves the "if" part of 3, and the "only if" part can be proved similarly.*